

W3D3
16/12/14

Random Spanning Trees

Geometry of Trees in the WSF

Remark from last time: $\tau_N = \sum_{n \leq N} \frac{1}{\chi_n} \chi(m, n)$ has infinite exp. $\rightarrow$
 $|E(X)| \rightarrow \infty$ requires transitivity.

Then μ^w -a.s. all trees T have the property: μ^w -a.s. the
 WSF on T is all of T .

In particular, if c is bounded, T is recurrent (Q5, ex. 2).

Observation If (T, c) is a transient network for a tree T , then
 $\exists e \in T$ s.t. both components of $T \setminus e$ are transient.

pf. Otherwise $\forall e \in T$, $T \setminus e$ has ≥ 1 recurrent comp. It follows
 that the tree has the form  ... but the line is also
 recurrent so the whole graph is.

pf. For any $e \in E$ define $A_e = \{e \text{ are both endpoints of } F\}$. Our goal -
 $\mu^w(e \in F, A_e) = 0$ [assume c bdd. $\Rightarrow$ this shows T is recurrent] by the obs.

Enumerate $G \setminus e$ by $e_1, e_2, \dots$ and let F_n be the σ -field
 $\sigma(\{f \in F\}, f = e_1, \dots, e_n)$. Let G_n be the graph where we contract
 $\{e_1, \dots, e_n\} \cap F$ and erase $\{e_1, \dots, e_n\} \cap F$.

By Kirchhoff $\mu^w(e \in F | F_n) = i(e, G_n) = c_e R(e \leftrightarrow e^+; G_n)$.

We claim that on A_e , this converges to 0 a.s., i.e.

$\mu^w(e \in F | F_n) \xrightarrow[n \rightarrow \infty]{a.s.} 0$. Given the claim, this limit is

~~$\mu^w(e \in F | F_\infty) \mathbb{1}_{A_e} = 0 \Rightarrow P(e \in F, A_e) = 0$~~

proof of the claim: if H is a transient network, $x \in V$,

and $\{e_1, \dots\} = E(H)$, $H_n = H / \{e_1, \dots, e_n\}$. If θ is a UCF

$x \rightarrow \infty$, $\forall \epsilon \exists N \sum_{i \geq N} \theta(e_i)^2 < \epsilon$. Project θ to $H_n \Rightarrow$

an $x \rightarrow \infty$ flow in H_n , $\epsilon(\theta, H_n) \leq \epsilon$. So $R(e \leftrightarrow \infty; G_n) \rightarrow 0$

on A_e , and similarly for e^+ (both are μ^w -a.s.).

By the triangle ineq., $R_{\text{eff}}(e^- \leftrightarrow e^+; G_n) \rightarrow 0$ on A_e .

An infinite simple path $v \rightarrow \infty$ is a ray. 2 rays are eq. if
 $|V_1 \Delta V_2| < \infty$. An "end" is an equivalence class.

Thm. On $\mathbb{Z}^d$ $d \geq 2$, each component of WSF has one end.

For a fin. planar graph $\exists$ bij. between spanning trees of G, G^* - $e \in T \Leftrightarrow e^* \in T^*$.

If $c(e)$ are conds. of G , put $c(e^*) = r(e)$

so that the weights $w(T) = \prod_{e \in T} c(e)$, $w(T^*) = \prod_{e^* \in T^*} r(e^*)$, $w(T) = w(T^*)$.

If G planar & infinite (locally finite $\rightarrow$ so is G^*), FSF on G is dual to WSF on G^* . On $\mathbb{Z}^2$ the dual is the graph itself, so $P(e \in T) = P(e^* \in T^*) = 1/2$.

Thm. If G is simple, planar, recurrent, G^* locally fin. & recurrent, the UST on G has 1 end.

pf. recurrent $\Rightarrow$ single tree in WSF/FSF. T has a bi-inf. path $\Rightarrow T^*$ has ≥ 2 components.

Lemma 1: Let $B(n)$ be a box of side length n in $\mathbb{Z}^d$, $d \geq 3$, then $\lim_{N \rightarrow \infty} \sup \{ R_{\text{eff}}(K \leftrightarrow \infty; \mathbb{Z}^d \setminus B(n)) : K \subseteq \mathbb{Z}^d \setminus B(n), |K| \leq N \} = 0$.

pf. We'll show $\lim_{N \rightarrow \infty} \sup \{ R_{\text{eff}}(K \leftrightarrow \infty; \mathbb{Z}^d) : |K| = N \} = 0$ later $N \times \mathbb{Z}^{d-1}$

$G = \mathbb{Z}^d$ is trans. Let θ be a UCF with fin. $\epsilon(\theta) = c$.

$\forall \epsilon \exists l \sum_{e \in E(G)} \theta^2(e) < \epsilon$. Let θ_x be the same flow $\left| \frac{d(e,0) \geq l}{d(e,0) \leq l} \right|$ $x \rightarrow \infty$, $K = \{x_1, \dots, x_N\}$.

$\exists$ at least $M = \frac{N}{(4l)^d}$ pts. in K of pairwise dist $\geq l$,

call them $x_1, \dots, x_M$. Take $\theta = \frac{1}{M} \sum_{i=1}^M \theta_{x_i}$. Then

$\epsilon(\theta) = \langle \theta, \theta \rangle = \frac{1}{M^2} \sum_{i,j} \langle \theta_{x_i}, \theta_{x_j} \rangle$. If $i \neq j$, $\langle \theta_{x_i}, \theta_{x_j} \rangle = \sum_{e, d(e, x_i) \leq l, d(e, x_j) \leq l} \theta_{x_i}(e) \theta_{x_j}(e) \leq 2\epsilon$. So, $\epsilon(\theta) = \frac{c}{M} + 2\epsilon \cdot c$.