

Alice sends measurement to Bob. If it's 1, Bob applies Z to y .

Julia Kempe: kempe@post.tau.ac.il

4 homework assignments.

Axioms of quantum:

- ① States - vectors in complex Hilbert space
- ② Dynamics - unitary matrices, $U \cdot U^\dagger = \mathbb{I}$ ($U^\dagger = U^\dagger$) [So the are real complex conjugate]
- ③ Measurements - "collapse" the state

$$\begin{aligned} \text{ket: } |0\rangle &= \begin{pmatrix} 1 \\ 0 \end{pmatrix} \\ \text{bra: } \langle 0| &= \begin{pmatrix} 1 & 0 \end{pmatrix}^\dagger = (1 \ 0) \end{aligned} \quad \begin{array}{l} \text{bra-ket} \\ \end{array} \quad \langle 0|0\rangle = (1 \ 0) \cdot \begin{pmatrix} 1 \\ 0 \end{pmatrix} = 1$$

$$|\psi\rangle = \alpha \cdot |0\rangle + \beta \cdot |1\rangle. \Rightarrow \text{Probability of seeing "0" is } |\alpha|^2 = |\langle 0|\psi\rangle|^2$$

We can choose in what basis to measure. For example,

$$\text{the } |\pm\rangle \text{ basis: } |\pm\rangle = \frac{1}{\sqrt{2}} (|0\rangle \pm |1\rangle).$$

In this basis, the probability of seeing "+" when measuring $|\psi\rangle$ is $|\langle +|\psi\rangle|^2$.

Say U is a unitary, such that $U|+\rangle = |0\rangle$, $U|-\rangle = |1\rangle$

$$|\langle +|\psi\rangle|^2 = |\langle +|U^\dagger U|\psi\rangle|^2 = |\langle 0|U|\psi\rangle|^2$$

$$[\text{In this case we have } U = H = \frac{1}{\sqrt{2}} \begin{pmatrix} 1 & 1 \\ 1 & -1 \end{pmatrix} = H^\dagger]$$

Partial measurements: measuring only a part of the qubits of $|\psi\rangle$

$$|\psi\rangle = \frac{1}{\sqrt{2}} (|00\rangle + |11\rangle). \quad "0" \rightarrow |00\rangle, \quad "1" \rightarrow |11\rangle$$

If we want to measure it over a different basis, say $|\pm\rangle$:

$$H \otimes \mathbb{I} |\psi\rangle = \frac{1}{2} (|00\rangle + |10\rangle + |01\rangle - |11\rangle)$$

$$\begin{array}{ll} "0" \rightarrow |0\rangle \frac{|0\rangle + |1\rangle}{\sqrt{2}} & "1" \rightarrow |1\rangle \frac{|0\rangle - |1\rangle}{\sqrt{2}} \end{array}$$

BB84 (Bennett - Brassard) [Quantum Cryptography]

Key distribution, for one-time pad.

① Alice sends randomly one of the following u to Bob: $|0\rangle, |1\rangle, |+\rangle, |-\rangle$

② Bob chooses a random basis from $\{|0\rangle, |1\rangle\}, \{|+\rangle, |-\rangle\}$.

* Notation: $|0\rangle \rightarrow \uparrow, |1\rangle \rightarrow \rightarrow, |+\rangle \rightarrow \searrow, |-\rangle \rightarrow \swarrow, |10\rangle, |1\rangle \rightarrow +, |1+\rangle \rightarrow X$

names	A	B	Result	Announce
"0"	$\uparrow$	+	0	+
"1"	$\rightarrow$	X	1	X
"0"	$\searrow$	+	1	X
"0"	$\uparrow$	X	0	+
"1"	$\rightarrow$	+	1	+
"0"	$\searrow$	+	1	X
"0"	$\rightarrow$	X	1	+
"1"	$\searrow$	X	1	X
"0"	$\uparrow$	+	0	+
"0"	$\rightarrow$	+	0	X

③ Bob measures the random bit in his basis

④ Alice publicly announces if the state of the bit she sent was in + or in X. Bob says ~~whether~~ whether he measured in the right basis.

⑤ Alice and Bob discard bits measured in a different basis.

- In expectation, they keep half the bits.

- If Eve measures the bits in the middle, with probability $\frac{1}{4}$ she will change Bob's result. To ensure they would catch her.

⑥ Alice and Bob verify consistency on half of the key.

* It is impossible to clone quantum states - No-Cloning Principle

Assume, we have a unitary U that can clone a state $|\psi\rangle$, for all $|\psi\rangle$

$$|\psi\rangle|0\rangle \xrightarrow{U} |\psi\rangle|\psi\rangle$$

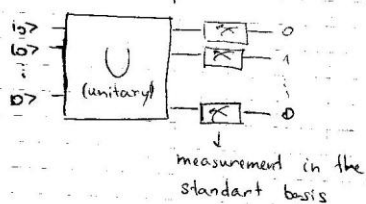
$$|0\rangle|0\rangle \xrightarrow{U} |0\rangle|0\rangle, |1\rangle|0\rangle \xrightarrow{U} |1\rangle|1\rangle$$

$$(\alpha|0\rangle + \beta|1\rangle)|0\rangle \xrightarrow{U} \text{want} \quad (\alpha|0\rangle + \beta|1\rangle) \otimes (\alpha|0\rangle + \beta|1\rangle)$$

$$\quad \quad \quad \text{get} \quad \quad \quad \neq \quad \alpha|00\rangle + \beta|11\rangle$$

- contradiction! $\square$

Quantum Computer



All quantum dynamics are unitary, and therefore reversible, whereas classical gates are not reversible. So we want to show that a quantum computer can simulate a classical computer.

We will first show that a classical computer can be simulated by a classical reversible computer.

A classical computer can be built only by 2 gates.

$$\text{NAND: } a, b \rightarrow 1 \oplus (a \cdot b)$$

$$\text{FANOUT: } a \rightarrow \begin{matrix} a \\ a \end{matrix}$$

$$\text{Toffoli gate (classical): } \begin{matrix} a & & a \\ b & \text{---} & b \\ c & \text{---} & c \oplus (a \cdot b) \end{matrix}$$

This gate is reversible (it is the inverse of itself).

We can use Toffoli gate to simulate both NAND and FANOUT

(and thus any classical computer) with only polynomial overhead!

$$\begin{matrix} a & & 1 \\ b & \text{---} & a \\ 1 & \text{---} & a \end{matrix} \quad 1 \oplus ab = \text{NAND}(a, b) \quad \begin{matrix} 1 & & 1 \\ a & \text{---} & a \\ 0 & \text{---} & a \end{matrix} = \text{FANOUT}(a)$$

The Toffoli gate is unitary, and we can therefore use it to simulate a classical computer with a quantum one.

Universality

(Classically: any function can be built from elementary gates)

- ① Any unitary can be built from CNOT and one-qubit unitaries.
- ② Any unitary can be approximated to any accuracy by $\{H, \text{CNOT}, \frac{\pi}{8}\text{-gate}\}$.

$$\frac{\pi}{8}\text{-gate: } \begin{pmatrix} e^{i\pi/8} & 0 \\ 0 & e^{-i\pi/8} \end{pmatrix} \quad (\text{or, up to phase } \begin{pmatrix} 1 & 0 \\ 0 & e^{-i\pi/4} \end{pmatrix}) \quad \left(\frac{\pi}{8}\right)^2 = \begin{pmatrix} 1 & 0 \\ 0 & -i \end{pmatrix}$$

$$\left(\frac{\pi}{8}\right)^4 = \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix} = 2\sigma_z$$

Def.: Norm of matrices: $\|U\| = \max_{\|v\|=1} \|Uv\|$ (for unitaries: always 1)

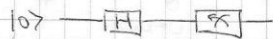
- When we say "approximate unitary U ", we mean we can build a unitary U_ϵ such that $\|U - U_\epsilon\| \leq \epsilon$, for any ϵ .

- There are many universal gate sets.

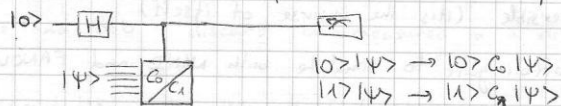
Theorem (Solovay-Kitaev): Any universal gate set can be efficiently approximated by any other.

Def.: BQP = the class of languages L that can be decided correctly with probability 75% by a polynomial size quantum computer.

We can simulate a random classical computer by:



* It is possible to postpone measurements: instead of measuring inside the computer and using a unitary G on C_1 according to the result we can use a special controlled unitary G_0/C_1 , and measure later.



Example: CNOT =

Black-box / Oracle:

classical: $x \mapsto f(x)$

classical reversible: $x \mapsto f(x) \oplus b$
(with Toffoli)

quantum: $|x\rangle \mapsto |f(x) \oplus b\rangle$